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## Datenstrukturen & Algorithmen

# Exercise Sheet 9 FS 16

**Exercise 9.1** *Minimum Spanning Trees.* 

Consider the following graph G = (V, E).



- a) Use Kruskal's algorithm to compute a minimum spanning tree for the graph below. Mark the edges contained in your solution.
- b) In a graph, a cycle is a walk starting and ending at the same vertex. A graph is *acyclic* if it contains no cycles. Show that every undirected acyclic graph G = (V, E) has at most |V| 1 edges.
- c) Show that the minimum spanning tree of a graph is uniquely determined if no edge weight occurs twice.

### **Exercise 9.2** Union-Find Structures.

In union-find structures, a set is represented by a tree. We consider the process of "unification by size" (see Chapter 6.2.2). We want to create a tree of height  $h \in \mathbb{N}$ . Describe how such a tree is generated with a sequence of UNION operations. How many UNION operations are necessary at least, and how many nodes does the resulting tree have?

### **Exercise 9.3** Longest increasing subsequence.

The following Java code contains an incomplete implementation of a dynamic program for the longest increasing subsequence problem. We are given a sequence  $A = A[0], \ldots, A[n-1]$  with pairwise different numbers  $A[i] \in \mathbb{N}^+$  for  $i = 0, \ldots, n-1$ . We compute a one-dimensional table T with n + 1 entries. The entry T[i] for i > 0 contains the smallest A[j] with  $j \in \{0, \ldots, n-1\}$ , such that A[j] is the last element of an increasing subsequence of length i. Initially, T[0] is set to  $-\infty$  and all other entries to  $\infty$ . In the *i*-th step, we search for the largest index j < i with T[j] < A[i]. We can use binary search to compute this index, because the entries of T are always monotonically increasing. If there is no such entry, we set j = 0. Then, we update the entry

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 $T[j+1] = \min(T[j+1], A[i])$ . After n steps, the length of the longest increasing subsequence is the largest index k with  $T[k] \neq \infty$ .

Complete the following Code snippet (lines 4, 5, 17, and 18).

```
1 static int binarySearch(int A[], int 1, int r, int key) {
\mathbf{2}
       while (1 < r) {
           int m = 1 + (r - 1 + 1)/2;
3
           if (A[m] >= key) _____
4
           else _
\mathbf{5}
       }
6
7
       return 1;
8
  }
9
10 static int computeTable(int A[], int size) {
                                         // DP-Tableau
11
       int[] T = new int[size+1];
       for (int i = 1; i <= size; i++) // Initialisierung</pre>
12
           T[i] = Integer.MAX_VALUE;
^{13}
14
       T[0] = Integer.MIN_VALUE;
       int 1 = 0;
15
16
       for (int i = ____; i ____; i = ____) {
17
           int j = binarySearch(___, ___, ___, ___
18
                                                      _);
           if ( A[i] < T[j+1] ) T[j+1] = A[i];</pre>
19
           if (l < j+1) l = j+1;
20
21
^{22}
       return 1;
23 }
```

### **Exercise 9.4** *Minimum spanning tree problem.*

In this exercise we are going to implement Kruskal's algorithm for the minimum spanning tree problem. This problem is defined as follows. Given an undirected graph G = (V, E) with the edge cost function  $c : E \to \mathbb{Q}^+$ , we search for an acyclic connected subset  $T \subseteq E$  with |T| = |V| - 1 (i.e., a tree (V,T)) whose total cost  $\sum_{e \in T} c(e)$  is minimum among all possible trees. The following image shows an example where the bold edges form a minimum spanning tree of the graph.

**Input** The first line contains only the number t of test instances. After that, we have exactly one line per test instance containing the description of the input graph G = (V, E) in the form  $n, m, u_1, v_1, c_1, \ldots, u_m, v_m, c_m$ . We have  $1 \le n, m \le 10000$  with  $V = \{1, \ldots, n\}$  and |E| = m. For every  $i, 1 \le i \le m$ , the numbers  $u_i, v_i \in \{1, \ldots, n\}$  define the edge  $\{u_i, v_i\} \in E$  having a cost of  $c_i, 1 \le c_i \le 1000$ .

**Output** For every test instance we output only one line. This line contains the cost of a minimum spanning tree.

#### Example

Input:

10 15 1 2 1 2 3 2 1 4 7 2	5 2 3 6 3 3 7 5 4 5 4 5 6 3 6 7	7 2 4 8 1 5 9 3 6 9 4 7	1068959103
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#### *Output:*

21

1

We provide you with a template. It contains the necessary code to read the input. Directions Implement the class UnionFind and Kruskal's algorithm. There is only one testset for 100 point in this exercise.

Hand-in: Wednesday, 4th May 2016 in your exercise group. 3