

**Datenstrukturen & Algorithmen****Exercise Sheet 2****FS 16****Exercise 2.1** *Estimating Asymptotic Running Time.*

Specify (as concisely as possible) the asymptotic running time of the following Java code fragments in  $\Theta$  notation depending on  $n \in \mathbb{N}$ .

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```
1 for (int i = 1; i <= n; i = i * 2) {
2     for (int j = n; j > 1; j -= 10)
3         ;
4 }
```

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1 for (int i = n; i > 0; i -= 5) {
2     for (int j = 0; j < i; j += 1) {
3         int k = 1;
4         while (k * k <= n)
5             k = k + 2;
6     }
7 }
```

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```
1 int f (int n) {
2     if (n == 1) return 1;
3     else {
4         for (int i = 1; i <= 2n; i++)
5             ;
6         return f(n/2)+1;
7     }
8 }
```

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**Exercise 2.2** *Recurrence Relations.*

Given a recurrence relation of the form

$$T(n) = \begin{cases} aT(\frac{n}{b}) + cn + d & \text{if } n > 1 \\ e & \text{if } n = 1 \end{cases}$$

with  $a, b, c, d, e \in \mathbb{N}$ ,  $a \neq b$ ,  $a \neq 1$  and  $b > 1$ . Find an expression with summations by telescoping and use the summation formula for geometric series ( $\sum_{i=0}^k q^i = \frac{q^{k+1}-1}{q-1}$  for  $q \neq 1$ ) to find a closed form. Prove your answer using mathematical induction. You can assume that  $n$  is a power of  $b$ .

*Please turn over.*

**Exercise 2.3** *Open Hashing.*

- a) Insert the keys 10, 18, 2, 4, 17 in this order into an initially empty hash table of size 7. Use open addressing with the hash function  $h(k) = k \bmod 7$  and resolve the conflicts using
- (i) linear probing (to the left)
  - (ii) quadratic probing (to the right, to the left, to the right, ...), and
  - (iii) double hashing with  $h'(k) = 1 + (k \bmod 5)$ .

For each method, provide the number of collisions. Which one is best for the above situation?

- b) Which problem occurs if the key 10 is removed from the hash tables in a), and how can you resolve it? Which problems occur if many keys are removed from a hash table?
- c) Provide a sequence of insert operations such that quadratic probing causes more collisions than linear probing. Specify a rule that prescribes how to form such a sequence of *arbitrary* length  $n$ . Use the hash function  $h(k) = k \bmod m$  for a prime number  $m$  with  $m \geq n$ .

**Exercise 2.4** *Cuckoo hashing.*

*Cuckoo hashing* is a hashing technique that guarantees constant time *in worst case* for both query and delete operations. The idea is to use two tables  $T_1$  and  $T_2$  of same size and two hash functions  $h_1$  and  $h_2$ . A new key  $x$  is inserted at position  $h_1(x)$  in  $T_1$ . In case of a collision, the previously stored key  $y$  is displaced to position  $h_2(y)$  in  $T_2$ . If this leads to another collision, the next key is again inserted at the appropriate position in  $T_1$ .

In some cases, this procedure continues forever, i.e. the same configuration appears after some steps of key displacements due to collisions. We will illustrate such a case in this exercise.

- a) Given two tables of size 5 each and two hash functions  $h_1 = k \bmod 5$  and  $h_2 = \lfloor k/5 \rfloor \bmod 5$ . Insert the following keys in the initially empty hash tables in this order: 3, 16, 18, 23, 1.
- b) Find another key, such that the insertion leads to an infinite sequence of key displacements. (In such a case, one would define two new hash functions, allocate two new tables and store the keys in these tables using the new hash functions.)

**Hand-in:** Wednesday, 9th March 2016 in your exercise group.