



## Data Structures & Algorithm

## Solutions to Sheet P5

## AS 16

### Solution P5.1 *Binary function value search.*

The solution was to use binary search for  $x$  in the range  $0 \dots x_{max}$ , comparing  $a$  to  $f(x)$  instead to `data[x]` as you would in a sorted array binary search. Otherwise the problem was essentially the same.

On the lecture website, you can find both non-recursive (`while`-loop based) and recursive solution, both running in time  $\mathcal{O}(n \log x_{max})$  and both implementing the algorithms from the lecture notes. Their source contains further comments on the implementation.

The time limit was set such that just computing all the values of  $f(x)$  for  $x = 0 \dots x_{max}$  even just once would take too long, so a binary search is necessary. Observe that since  $\log_2(x_{max}) \simeq 24.25$ , less than 30 calls to  $f$  are necessary for any input, and contrast that to  $x_{max} = 2 \cdot 10^7$  calls required to a linear scan of the range.

### Data

`judge1`  $n = 11$  values including 0 and 1 (the lowest possible values),  $500000000 = f(x_{max})$  and  $500000001 = f(x_{max}) + 1$  and some random values in between.

**Notes on submitted solutions.** The most common problem was the input value  $f(x_{max})$ . Some solutions maintained a search range with variables such as `min` and `max`, but only worked for values `min .. max - 1`. While it can be natural to work with these conditions, the solutions also set the starting search range to `min = 0` and `max = x_max`, missing the maximal value.

A second problem was that some solutions contained a loop like `while (min != max) {...}`, but then they would also somehow set `max = min - 1` during the search, e.g. via `max = (max + min) / 2 - 1` after having `max == min + 1`. This would result in an infinite loop. One valid solution is to replace `(min != max)` with `(min < max)`. Note that infinite recursion fails with `RUNTIME ERROR` in the judge.

**Notes on the function  $f$ .** (Extra material, not necessary to solve the problem)

The function is actually  $f(x) = xl(x)$ , where  $l(x) = \lfloor \log_2 x \rfloor + 1$ , which is the number of binary digits of  $x$ . Knowing this was in no way important to solve the problem, but it is theoretically possible to compute the inverse of  $f$  almost directly. (Again, you are encouraged to try.) However, it is certainly not easier than the binary search program and would be significantly faster only for very large  $x_{max}$ .